

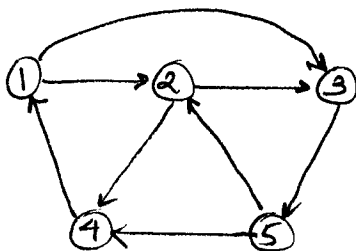
## SOLUTIONS FOR HW # 7

① Given a directed graph  $G(V, E)$ , find a minimum number of vertices, such that eliminating them from  $G$ , makes the graph cycle-free.

② for  $\forall v \in V$  ( $i$ , from 0 to  $|V|$ )  
for each vertex set  $S$  with size  $i$  ( $|S| = i$ )  
remove  $S$  from  $G$  and get  $G'$   
DFS( $G'$ )  
if  $G'$  has no cycle  
minimum number  $\leftarrow i$   
FVS  $\leftarrow S$   
stop.

Complexity:  $\sum_{i=0}^{|V|} \binom{|V|}{i} O(|V| + |E|) = 2^{|V|} O(|V| + |E|)$  Exponential.

③



$i=0, 1$   $G'$  has a cycle.

$i=2$   $S = \{1, 2\}$   $G'$  has no cycle.

So the minimum number = 2

FVS:  $\{1, 2\}$

!!! by can verify for all other subsets  
(by inspection) and we conclude that

this is the optimum.

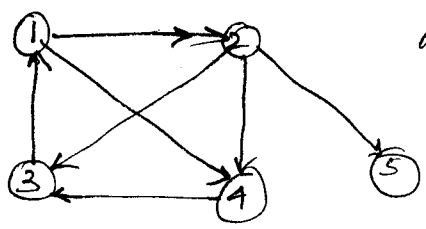
④ Heuristic Algorithm,

1. Select a vertex with the highest degree in  $G$ .
2. Remove this from  $G$ .
3. Apply DFS( $G$ )
4. (If  $G$  has no cycle). stop else repeat 2 and 3.

Each iteration has atmost  $|V|$  complexity

selecting vertex:  $O(|V|)$       DFS =  $O(|V| + |E|)$   
 total complexity =  $O(|V|) * (O(|V|) + O(|V| + |E|))$   
 =  $O(|V| * (|V| + |E|))$ .

⑤



apply heuristic

select ②  
 then select ①

$FVS = \{1, 2\}$       size = 2

But had you selected vertex 1 (size = 1)

It would have been the optimum.

Optimum  $FVS = \{1\}$

⑥

program & runtime results.